

# Finding the Seed of Uniform Attachment Trees

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July 26, 2019

## Network Archaeology on Random Trees

Setup

Results

Skecth of the proofs

# Introduction

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This problem was popularized by Shah-Zamah [6].

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- ▶  $T_\ell = S_\ell$  ;
- ▶  $T_i$  is obtained by **joining vertex  $i$  to a vertex of  $T_{i-1}$  chosen uniformly at random**, independently of the past.

## Influence of the seed

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Curien, Duquesne, Kortchemski and Manolescu: YES. [4]

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Given  $T_n(P_\ell)$ ,  $n \gg 1$  we can find a set

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$$\mathbb{P} \{H_n \subset P_\ell\} \geq 1 - \epsilon .$$

## Theorem 2: $S_\ell = E_\ell$

For  $\ell \geq \max \left\{ C, \frac{8}{\gamma} \right\} \log \frac{1}{\epsilon}$  we have the following:

Given  $T_n(E_\ell)$ ,  $n \gg 1$  we can find a set

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### Theorem 3: $S_\ell = T_\ell$

There exist  $c_1$  and  $c_2$  such that the following holds. Let

$$\ell \geq c_1 \log^2 \frac{1}{\epsilon}.$$

Given  $T_n(T_\ell)$ ,  $n \gg 1$  we can find a set

$H_n = H_n(T_n, \epsilon) \subset \{1, \dots, n\}$  with  $|H_n| \geq \ell / [c_2 \log(\ell/\epsilon)]$  such that

$$\mathbb{P}\{H_n \subset T_\ell\} \geq 1 - \epsilon.$$

## Theorem 4

### Theorem

Let  $\epsilon \in (0, e^{-e^2})$ . Suppose that  $T_n$  is a uniform attachment tree with seed  $S_\ell = P_\ell$  or  $S_\ell = E_\ell$  for  $\ell \leq \frac{\log(1/\epsilon)}{\log \log(1/\epsilon)}$ . Then, for all  $n \geq 2\ell$ , any seed-finding algorithm that outputs a vertex set  $H_n$  of size  $\ell$  has

$$\mathbb{P} \left\{ |H_n \cap S_\ell| \leq \frac{\ell}{2} \right\} \geq \epsilon .$$

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Let us define what means be more central.

# Rooted tree and Induced subtree

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A **rooted tree**  $(T, v)$  is the tree  $T$  with a distinguished vertex  $v \in V(T)$ .

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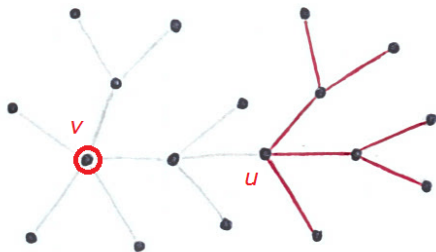
A **rooted tree**  $(T, v)$  is the tree  $T$  with a distinguished vertex  $v \in V(T)$ .

The **subtree induced by  $u$**   $(T, v)_{u\downarrow}$  is the subtree of  $(T, v)$  which grows from  $u$  in the opposite direction of  $v$ .

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## Centrality: Definition

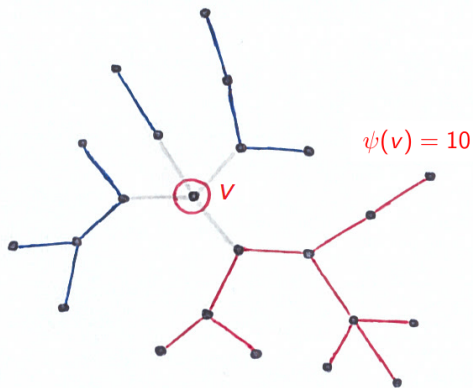
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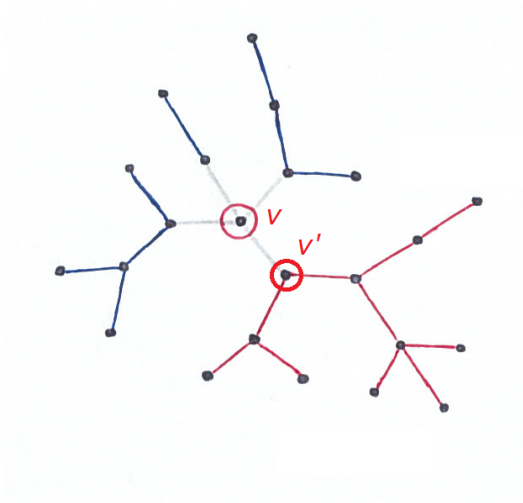




# Centrality

Given  $v$ , we denote  $v'$  to be some vertex in  $N(v)$  such that

$$\psi(v) = |(T, v)_{v' \downarrow}| .$$



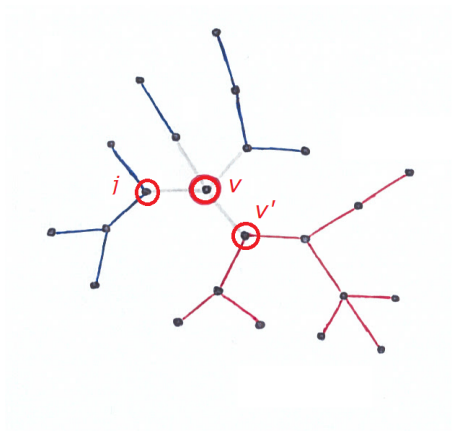
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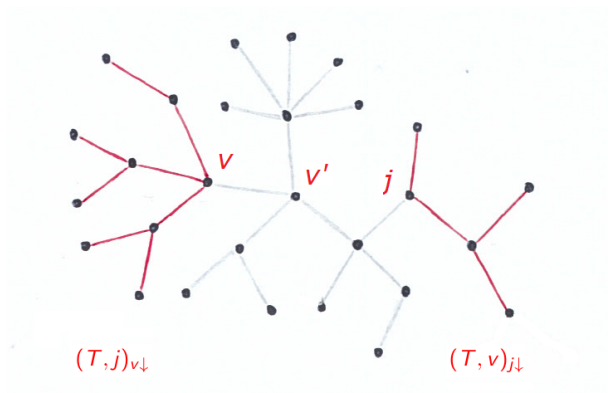
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Case 2: When  $v'$  is between  $v$  and  $j$  we have  $\psi(v) \leq \psi(j)$  if

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- ▶  $|(T, j)_{v\downarrow}| \geq |(T, v)_{j\downarrow}|;$



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More precisely

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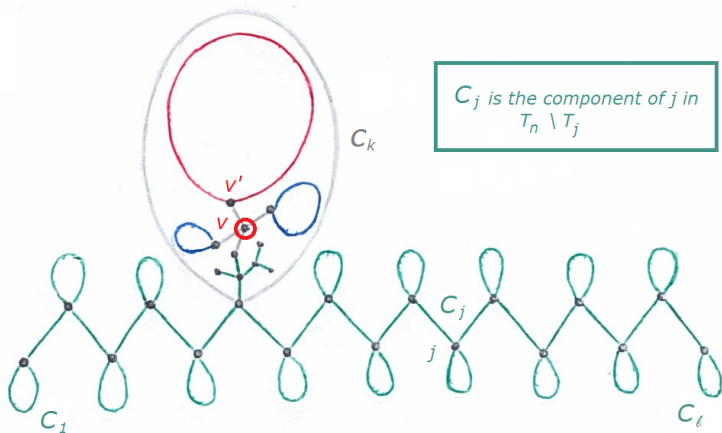
Let us prove that the complement has small probability.

## Sketch of the proof

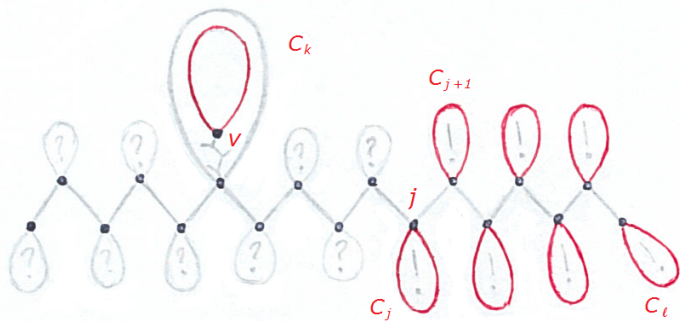
By Union Bound we have

$$\begin{aligned} \mathbb{P} \left\{ \min_{\ell < i \leq n} \psi(i) \leq \max_{\ell\gamma/2 \leq j \leq \ell(1-\gamma/2)} \psi(j) \right\} \\ \leq \sum_{j=\gamma\ell/2}^{(1-\gamma/2)\ell} \mathbb{P} \left\{ \min_{\ell < i \leq n} \psi(i) \leq \psi(j) \right\} \\ \leq \sum_{j=\gamma\ell/2}^{(1-\gamma/2)\ell} \sum_{k=1}^{\ell} \mathbb{P} \left\{ \exists v \in C_k \setminus \{k\} : \psi(v) \leq \psi(j) \right\}. \end{aligned}$$

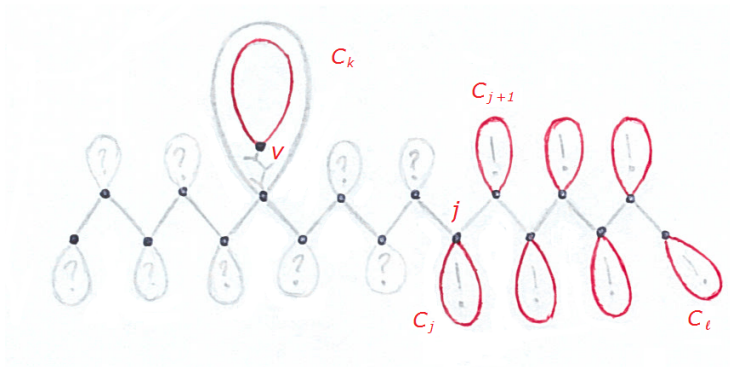
First case:  $(v', v, j)$







Second case:  $(v, v', j)$



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The value of  $l$  arise from a optimization of the bounds.

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Thank  
you

